

Module 15: Screening

Information Economics (Ec 515) · George Georgiadis

Non-Linear Pricing

- Quasi-linear model
- Consumer (agent)
 - Type: θ (taste)
 - Utility $u = \theta q - t$ (taste * quality - price)
- Firm's Profit: $t - c(q)$ ($c', c'' > 0$)

Two approaches:

1. Revelation mechanism:

- Firm asks agent for type θ .
- Consumer reports $\tilde{\theta}$ (not necessarily $\tilde{\theta} = \theta$).
- Mechanism specifies $q(\tilde{\theta}), t(\tilde{\theta})$.
- *Objective*: Design mechanism such that the consumer finds it optimal to report $\tilde{\theta} = \theta$.

(a) Taxation mechanism:

- Firm offers menu of contracts $T(q)$.
 - Consumer, knowing his type θ , picks most favorable contract.
 - *Objective*: Design menu to maximize expected profit.
- **Lemma 1** (*Revelation principle & Taxation Principle*): The two approaches are equivalent.

Proof.

Revelation principle \Rightarrow *Taxation principle*:

- Suppose the revelation mechanism $q(\cdot)$ and $t(\cdot)$ is incentive compatible (*i.e.*, type θ finds it optimal to report $\tilde{\theta} = \theta$).

- Let

$$T(q) = \begin{cases} t(\theta) & \text{if } q = q(\theta) \\ \infty & \text{otherwise} \end{cases}$$

- Since $q(\cdot)$ and $t(\cdot)$ is incentive compatible, type θ buys $q(\theta)$.

Taxation principle \Rightarrow *Revelation principle*:

- Fix some menu $T(q)$. Now let

$$q(\theta) = \max_q \{\theta q - T(q)\} \quad \text{and} \quad t(\theta) = T(q(\theta)) .$$

- Type θ finds it optimal to report

$$\tilde{\theta} = \arg \max_{\theta'} \{\theta q(\theta') - t(\theta')\} = \arg \max_{\theta'} \{\theta q(\theta') - T(q(\theta'))\} = \theta$$

□

- The revelation mechanism has proved to be more convenient to use!

- Utility of a type θ who reports $\tilde{\theta}$:

$$\begin{aligned} u(\theta; \tilde{\theta}) &= \theta q(\tilde{\theta}) - t(\tilde{\theta}) \\ u(\theta) &= u(\theta; \theta) \end{aligned}$$

- Principal's problem: Choose $q(\cdot)$ and $t(\cdot)$ to maximize

$$\mathbb{E}_\theta [t(\theta) - c(q(\theta))]$$

subject to

$$\begin{aligned} u(\theta; \theta) &\geq 0 && \text{(IR)} \\ u(\theta; \theta) &\geq u(\theta; \tilde{\theta}) \quad \forall \tilde{\theta} && \text{(IC)} \end{aligned}$$

First-best Outcome:

- Suppose that the principal knows the agent's type.
- Consumer gets 0 utility.
 - So $t(\theta) = \theta q(\theta)$, so the principal solves $\max_{q(\cdot)} \mathbb{E}_\theta [\theta q(\theta) - c(q(\theta))]$.
 - Pointwise maximization w.r.t q yields: $\theta = c'(q(\theta))$.
- Does this satisfy (IC) ?
 - No. Suppose $\theta \in \{\theta_L, \theta_H\}$.
 - Then θ_i receives q_i and assuming $q_L > 0$, we have

$$\begin{aligned} u(\theta_H; \theta_L) &= \theta_H q_L - t_L = (\theta_H - \theta_L) q_L + \theta_L q_L - t_L \\ &= \underbrace{(\theta_H - \theta_L) q_L}_{>0} + \underbrace{u(\theta_L; \theta_L)}_{=0} \\ &> 0 = u(\theta_H; \theta_H) \end{aligned}$$

- *i.e.*, type θ_H wants to imitate type θ_L .

Screening with Two Types

- Suppose $\theta \in \{\theta_L, \theta_H\}$, where $\theta_H > \theta_L$
- $\Pr\{\theta = \theta_H\} = 1 - \pi$
- Principal maximizes

$$\max_{q_H, t_H, q_L, t_L} \{(1 - \pi) [t_H - c(q_H)] + \pi [t_L - c(q_L)]\}$$

subject to

$$\begin{aligned} \theta_H q_H - t_H &\geq 0 && (IR_H) \\ \theta_L q_L - t_L &\geq 0 && (IR_L) \\ \theta_H q_H - t_H &\geq \theta_H q_L - t_L && (IC_H) \\ \theta_L q_L - t_L &\geq \theta_L q_H - t_H && (IC_L) \end{aligned}$$

- Assume $q_L > 0$. (Otherwise, just serve θ_H efficiently and extract all rents.)

○ *First-best Outcome:*

- IR_H and IR_L bind: $t_H = \theta_H q_H$ and $t_L = \theta_L q_L$.
- Objective function becomes: $\max_{q_H, q_L} \{(1 - \pi) [\theta_H q_H - c(q_H)] + \pi [\theta_L q_L - c(q_L)]\}$
- First order conditions: $c'(q_H^{fb}) = \theta_H$ and $c'(q_L^{fb}) = \theta_L$; *i.e.*, marginal cost = marginal benefit.

○ *Proposition:* (IR_L) and (IC_H) are binding, while the other two constraints can be replaced by the monotonicity constraint $q_H \geq q_L$.

Proof.

○ (IR_H) is slack:

$$\begin{aligned} \theta_H q_H - t_H &\geq \theta_H q_L - t_L \text{ (by } IC_H) \\ &> \theta_L q_L - t_L \\ &\geq 0 \text{ (by } IR_L) \end{aligned}$$

- (IR_L) binds: If not, increase t_L until it binds to increase the principal's profit.
- Monotonicity: Adding (IC_H) and (IC_L) yields

$$\begin{aligned} \theta_H q_H - \theta_L q_H &\geq \theta_H q_L - \theta_L q_L \\ (\theta_H - \theta_L) q_H &\geq (\theta_H - \theta_L) q_L \\ q_H &\geq q_L \end{aligned}$$

- (IC_H) binds: If not, increase t_H by ϵ because (IR_H) is slack.
- (IC_L) is redundant:

$$\begin{aligned} t_H - t_L &= \theta_H (q_H - q_L) \text{ (by } IC_H) \\ &\geq \theta_L (q_H - q_L) \text{ (by monotonicity)} \end{aligned}$$

□

○ Program becomes:

$$\max_{q_H, t_H, q_L, t_L} \{(1 - \pi) [t_H - c(q_H)] + \pi [t_L - c(q_L)]\}$$

subject to

$$\theta_L q_L - t_L = 0 \quad (IR_L)$$

$$\theta_H q_H - t_H = \theta_H q_L - t_L \quad (IC_H)$$

$$q_H \geq q_L \quad (\text{monotonicity})$$

- Relax program by ignoring (monotonicity), and using (IR_L) and (IC_H) :

$$\begin{aligned} & (1 - \pi) [\theta_H q_H - (\theta_H - \theta_L) q_L - c(q_H)] + \pi [\theta_L q_L - c(q_L)] \\ = & (1 - \pi) \left[\underbrace{\theta_H q_H - c(q_H)}_{\text{welfare}} \right] + \pi \left[\underbrace{\theta_L q_L - c(q_L)}_{\text{welfare}} - \underbrace{\frac{1 - \pi}{\pi} (\theta_H - \theta_L) q_L}_{\text{expected info rents}} \right] \end{aligned}$$

- *First-order conditions:* $c'(q_H) = \theta_H$ and $c'(q_L) = \theta_L - \frac{1 - \pi}{\pi} (\theta_H - \theta_L)$

- Observe that $q_H = q_H^{fb}$.

- But $c'(q_L) < \theta_L$, which implies that $q_L < q_L^{fb}$.

- Check monotonicity:

$$\begin{aligned} c'(q_H) &= \theta_H \\ &> \theta_L - \frac{1 - \pi}{\pi} (\theta_H - \theta_L) = c'(q_L) \\ \implies q_H &> q_L \end{aligned}$$

- Once we have determined q_H and q_L , we can back out the transfers from (IR_L) and (IC_H) :

- $t_L = \theta_L q_L$

- $t_H = t_L + \theta_H (q_H - q_L)$

Properties:

- Low type is inefficiently underserved.

- Makes it less attractive for the high type to imitate the low type; *i.e.*, lowers rents of high type.
- *e.g.*, squeeze passengers in economy class.
- Lowest type gets no surplus (that would be a waste).
- Efficiency at the top.
 - Low type cannot “afford” to mimic the high type.
 - Serve him optimally (and tax him).
- High type indifferent between contracts.
 - Ensured by making economy class uncomfortable.
- Quality / quantity increases in type.

Screening with a Continuum of Types

- $\theta \sim f(\cdot)$ with support on $[\underline{\theta}, \bar{\theta}]$.
 - The agent’s “type” θ can now take a continuum of values (instead of 2).
- *Recall:* Utility of a type θ who reports $\tilde{\theta}$:

$$\begin{aligned} u(\theta; \tilde{\theta}) &= \theta q(\tilde{\theta}) - t(\tilde{\theta}) \\ u(\theta) &= u(\theta; \theta) \end{aligned}$$

- Principal’s problem:

$$\max_{q(\theta), t(\theta)} \mathbb{E}_{\theta} [t(\theta) - c(q(\theta))]$$

subject to

$$\begin{aligned} u(\theta; \theta) &\geq 0 && \text{(IR)} \\ u(\theta; \theta) &\geq u(\theta; \tilde{\theta}) \quad \forall \tilde{\theta} && \text{(IC)} \end{aligned}$$

Theorem. $q(\cdot), t(\cdot)$ is incentive compatible if and only if

$$u(\theta) = u(\underline{\theta}) + \int_{\underline{\theta}}^{\theta} q(s) ds \quad (\text{Payoff Equivalence})$$

$$q(\theta) \text{ is increasing} \quad (\text{Monotonicity})$$

Proof.

◦ Fix $\theta' > \theta$.

◦ Only if:

– (IC) implies that

$$\left. \begin{array}{l} \theta'q(\theta') - t(\theta') \geq \theta'q(\theta) - t(\theta) \\ \theta q(\theta) - t(\theta) \geq \theta q(\theta') - t(\theta') \end{array} \right\} \implies \theta' [q(\theta') - q(\theta)] \geq t(\theta') - t(\theta) \geq \theta [q(\theta') - q(\theta)]$$

and hence $q(\theta') - q(\theta) \geq 0$.

– Payoff equivalence follows from the Envelope Theorem:

$$\frac{du}{d\theta}(\theta) = \frac{\partial u}{\partial \theta}(\theta; \theta) + \underbrace{\frac{\partial u}{\partial \tilde{\theta}}(\theta; \tilde{\theta}) \Big|_{\tilde{\theta}=\theta}}_{=0 \text{ (FOC at } \tilde{\theta}=\theta)} = q(\theta)$$

(Real proof in Milgrom and Segal (ECTA, 2002).)

◦ If:

$$\begin{aligned} u(\theta') &= u(\theta) + \int_{\theta}^{\theta'} q(s) ds \\ &\geq u(\theta) + \int_{\theta}^{\theta'} q(\theta) ds \quad (\text{by monotonicity}) \\ &= u(\theta) + (\theta' - \theta) q(\theta) \\ &= u(\theta'; \theta) \end{aligned}$$

□

◦ Principal's Problem:

$$\max_{q(\theta), t(\theta)} \mathbb{E}_{\theta} [t(\theta) - c(q(\theta))]$$

subject to

$$u(\theta) = u(\underline{\theta}) + \int_{\underline{\theta}}^{\theta} q(s) ds \quad (\text{Payoff Equivalence})$$

$$u(\underline{\theta}) = 0 \quad (\text{IR low})$$

$$q(\theta) \text{ is increasing} \quad (\text{Monotonicity})$$

- $u(\theta) = \theta q(\theta) - t(\theta)$ implies that $t(\theta) = \theta q(\theta) - u(\theta)$ so that the objective function can be re-written as $\mathbb{E}[\theta q(\theta) - c(q(\theta)) - u(\theta)]$

- Notice:

$$\begin{aligned} \mathbb{E}[u(\theta)] &= \int_{\underline{\theta}}^{\bar{\theta}} u(\theta) dF(\theta) \\ &= - \int_{\underline{\theta}}^{\bar{\theta}} u(\theta) [1 - F(\theta)]' d\theta \\ &= - \underbrace{[u(\theta) \bar{F}(\theta)]_{\underline{\theta}}^{\bar{\theta}}}_{=0} + \int_{\underline{\theta}}^{\bar{\theta}} \underbrace{u'(\theta)}_{=q(\theta)} \bar{F}(\theta) d\theta \\ &= \mathbb{E} \left[q(\theta) \frac{1 - F(\theta)}{f(\theta)} \right] \end{aligned}$$

- We can write the principal's profit as

$$\mathbb{E} \left[q(\theta) \left[\theta - \frac{1 - F(\theta)}{f(\theta)} \right] - c(q(\theta)) \right]$$

- Ignore the monotonicity constraint (for now), and maximize this pointwise with respect to q .

$$- \text{First-order condition: } \underbrace{\theta - \frac{1 - F(\theta)}{f(\theta)}}_{MR(\theta)} - \underbrace{c'(q(\theta))}_{MC(\theta)} = 0$$

- As long as $MR(\theta)$ is increasing, $q(\theta)$ is increasing in θ , and the monotonicity constraint is satisfied.

- *e.g.*, true if θ follows a uniform or exponential distribution.

- Selling $q(\theta)$ generates:

- Surplus $\theta q(\theta)$

- Costs $c(q(\theta))$
- Consumer rents $\frac{1-F(\theta)}{f(\theta)}q(\theta)$
- How to back out transfers $t(\cdot)$?
 - The above maximization problem yields $q(\theta)$.
 - We can then compute $u(\theta) = \int_{\underline{\theta}}^{\theta} q(s) ds$.
 - Finally, we obtain $t(\theta) = \theta q(\theta) - u(\theta)$.

Example 1: Quadratic Costs

- Setup: $\theta \sim U[0, 1]$; $c(q) = \frac{q^2}{2}$
- Marginal Revenue: $MR(\theta) = 2\theta - 1$
- Marginal Cost: $c'(q) = q$
- Optimal Contract: $q(\theta) = [2\theta - 1]^+$
 - Then $u(\theta) = \int_{\frac{1}{2}}^{\theta} (2s - 1) ds = \frac{1}{4} - \theta(1 - \theta)$; and
 - Transfers $t(\theta) = \theta(2\theta - 1) - \frac{1}{4} + \theta(1 - \theta) = \theta^2 - \frac{1}{4}$.

Example 2: Linear Costs

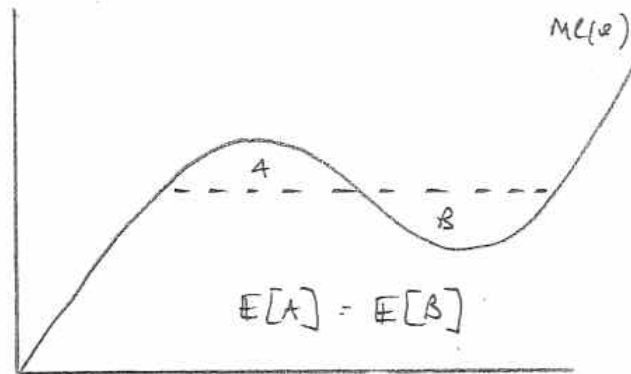
- Setup: $c(q) = cq$ with $q \in [0, 1]$
- No haggling Theorem:

$$q(\theta) = \begin{cases} 1 & \text{if } MR(\theta) \geq c \\ 0 & \text{otherwise} \end{cases}$$

Further Remarks:

- One agent or many agents
 - It doesn't matter.
 - If types are independent, then principal cannot gain by linking games.
 - If types are correlated, then principal can fully extract surplus by linking games (Cremer and Maclean, 1985).
- Generalizes to n agents with independent information: Optimal auctions

- Quality or Quantity:
 - With one consumer, q could also be quantity.
 - With many consumers, games are linked if costs are convex, because $c(\sum_i q_i)$ (Segal, AER 2003).
- What if marginal revenue $MR(\theta)$ is not increasing?
 - Need to “iron” it.
 - Two approaches:
 - * Optimal control (Guesnerie and Laffont, 1984)
 - * Convex hull (Myerson, 1981)
 - Convex hull approach: Replace $\int^\theta MR(s) ds$ with smallest convex envelope (draw picture).



- * x-axis: θ
- * y-axis: $MR(\theta)$
- * We choose a value to flatten the allocation function $q(\theta)$ such that $\mathbb{E}[A] = \mathbb{E}[B]$.

References

Board S., (2011), Lecture Notes.

Bolton and Dewatripont, (2005), *Contract Theory*, MIT Press.