

A Robust Folk Theorem for the Prisoner's Dilemma*

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We prove the folk theorem for the Prisoner's dilemma using strategies that are robust to private monitoring. From this follows a limit folk theorem: when players are patient and monitoring is sufficiently accurate, (but private and possibly independent) any feasible individually rational payoff can be obtained in sequential equilibrium. The strategies used can be implemented by finite (randomizing) automata. *Journal of Economic Literature* Classification Numbers C72, C73, D82.

1. INTRODUCTION

The folk theorem for discounted repeated games states that every payoff vector that is feasible and individually rational is an equilibrium payoff when players are sufficiently patient. A proof of the folk theorem first appeared in Fudenberg and Maskin [7] for subgame-perfect equilibria of repeated games with *perfect monitoring*. Perfect monitoring means that the history of chosen actions is always common knowledge among the players. In many important economic applications, players monitor one another imperfectly: each observes a noisy signal of the actions chosen by others. Thus, starting with the early papers by Radner [16] and Green and Porter [9], attention turned to repeated games with imperfect monitoring. The

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early applications were to situations in which monitoring, while imperfect, was public: the random payoff-relevant outcome in each stage was assumed common knowledge among the players. Exploiting the structure of public monitoring, dynamic programming techniques can be used to characterize the set of perfect public equilibrium payoffs as in Abreu, Pearce, and Stacchetti [3] and [4]. This approach to imperfect monitoring culminated in the folk theorem of Fudenberg, Levine, and Maskin [5] which identified conditions on the public monitoring technology which ensured that all feasible and individually rational payoffs could be supported in equilibrium.

In repeated games with monitoring by *privately* observed signals, these techniques do not apply and whether the folk theorem extends is still an open question. Indeed, for some games, whether there are *any* sequential equilibria different from repetition of stage-game Nash profiles is unresolved.¹ Numerous negative results emphasize the difficulties involved. Matsushima [14] considers repeated play of stage games with a unique Nash equilibrium and monitoring by *conditionally independent* private signals. Conditional independence means that for each action profile, the players' private signals are distributed independently of one another. Each individual's private signal may be arbitrarily informative about the realized action profile, but conditional on that action profile, each player's signal is uninformative about the signals observed by other players. Matsushima shows that if there is any pure-strategy equilibrium different from repetition of the stage game Nash profile, it must involve conditioning on payoff-irrelevant history. Specifically, there must be a player i and a pair of histories for i which give rise to identical beliefs over the opponents' continuation strategies but which nevertheless induce distinct continuation play by i . In particular, this implies that equilibria cannot be "strict."²

Many of the strategies used to prove folk theorems in environments with public monitoring fail to be even approximate equilibria when monitoring is at all imperfect but conditionally independent. For example, consider the strategies used by Fudenberg and Maskin [8] to prove the folk theorem under perfect monitoring. These strategies begin in a cooperative phase in which players play a deterministic sequence of pure action profiles until some player deviates from that sequence. Such a deviation triggers a punishment in which the deviating player is minmaxed. Continuation strategies are constructed so that the non-deviating players have an in-

¹For a particularly problematic example, see the discussion in Mailath and Morris [13] of the "convention game" originally studied by Shin and Williamson [18]

²Bhaskar [1] derives an even stronger necessary condition from a model in which each player's payoffs are randomly perturbed in each period independently of history and are private information to that player. In a repeated game with overlapping generations of players (hence private monitoring), his condition implies that any equilibrium must be a repetition of stage-game Nash outcomes.

centive to carry out the punishment and return to the cooperative phase. When the discount factor is close enough to one, each player has a strict incentive to follow his equilibrium strategy after every history.

It is this strictness property that implies that even the slightest conditionally independent private monitoring imperfections destroy the equilibrium. To see this, consider a player i who has detected a deviation by player j in stage 1. Since i knows that j is following his equilibrium strategy, i knows that j has not deviated and in fact that no player has deviated. Since the equilibrium was constructed so that i has a strict incentive to cooperate after histories in which no player has deviated, i will not punish but will instead continue to cooperate. But this unwillingness to punish signals of bad behavior eliminates the necessary incentives to cooperate in the first place.

This naturally raises the question of what equilibria of games with perfect monitoring are *robust* to private monitoring imperfections. That is, for which (if any) equilibria is there guaranteed to be a nearby equilibrium when monitoring is nearly perfect, but private. In this paper, we focus on the repeated prisoners' dilemma and prove a robust folk theorem: for sufficiently patient players, every feasible, individually rational payoff can be achieved by equilibrium strategies that are robust to private monitoring. This implies a limit folk-theorem: For monitoring technologies sufficiently accurate and discount factors close enough to one, any individually rational payoff can be achieved in a sequential equilibrium of the private monitoring game.

Related results have been obtained elsewhere. Sekiguchi [17] was the first to show that the mutual cooperation payoff can be achieved in a class of prisoners' dilemma games when monitoring is nearly perfect.³ Mailath and Morris [13] show that some trigger strategies are robust to monitoring imperfections that are approximately "public." Monitoring is approximately public when a player's realized signal is informative of the signals observed by others (even after conditioning on the realized action.) Conditional independence is therefore ruled out.

To prove that trigger-strategy and other strict equilibria are robust, the problem of discontinuous conditional beliefs must be circumvented. Conditional beliefs about opponents' continuation strategies after histories involving a deviation change discontinuously when slight monitoring imperfections are introduced. This can be seen in the above discussion of the Fudenberg and Maskin [8] strategies. When monitoring is perfect and player i observes a deviation in stage 1, he is certain that his opponents will move into the punishment phase. But when monitoring is conditionally indepen-

³Bhaskar and Obara [2] sharpen the analysis and show that the payoff restriction used in Sekiguchi [17] is not necessary. When public correlation devices are available, these strategies can also be used to prove the folk theorem.

dent, following the corresponding (private) history player i is nearly certain that his opponents will continue to cooperate. This observation, together with the Matsushima result makes it clear that strict equilibria cannot be robust to *all* private monitoring imperfections.

Mailath and Morris [13] restrict attention to monitoring technologies in which there is sufficient correlation in the players' private signals, i.e. nearly public monitoring. When monitoring is approximately public, player i will be nearly certain that the opponents have also observed a deviation and will begin punishing, even though it is common knowledge among the players that no deviation has occurred. Notice that conditional beliefs are no longer discontinuous once the monitoring imperfections are restricted in this way.⁴

Our approach to the folk theorem deals with discontinuous conditional beliefs in a simpler way. We construct stationary strategies which have the property that each player is indifferent among each of his actions no matter what private history his opponent has observed. It is then irrelevant how conditional beliefs are altered by the monitoring imperfections since the players' continuation strategies will be a best-reply to *every* conditional belief. This in turn implies that our strategies will be robust to all monitoring imperfections, including conditional independence and even negative correlation. We show that a large set of payoffs, including mutual cooperation, can be supported by such equilibria. However, not all individually rational payoffs can be supported in this way. To obtain the full folk theorem set of payoffs, we show how these robust equilibria generate continuation values that can enforce behavior yielding payoffs outside the original set for a sufficiently long but finite length of time.

A recent paper by Piccione [15] employs a technique similar to ours. In the repeated prisoners' dilemma Piccione uses dynamic programming techniques over the infinite state-space of private histories to construct the mixed strategies necessary to maintain indifference. His strategies can be used to approximate most of the feasible, individually rational payoff set. Our approach uses stationary behavior strategies that condition only on one period of history. This dramatically simplifies the dynamic program yielding a system of four equations which can easily be solved for equilibrium mixtures.⁵ Beyond the computational simplicity, the simple structure makes this a more promising direction for results in more general games. We demonstrate the flexibility of our approach by analyzing some more general games in section 5.

⁴There is still an additional complication to overcome in establishing that a strict equilibrium is robust to approximate public monitoring. Conditional beliefs must move continuously *uniformly* across the infinite set of histories.

⁵We also obtain most of the folk theorem payoffs exactly.

In section 2, the perfect monitoring folk theorem is proven for a large subset of the feasible payoff set using simple two state mixed strategies. In section 3, it is shown that the strategies used are robust, and the limit folk theorem is established for these payoffs. In section 4 we show how to augment these strategies to obtain the full set of feasible individually rational payoffs. In section 5, we consider more general games. For 2 player games, we provide a sufficient condition on stage payoffs under which our approach will apply. We also analyze a symmetric N -player prisoners' dilemma and show how to extend our approach and obtain nearly cooperative payoffs for high discount factors and sufficiently accurate monitoring. Finally, section 6 concludes.

2. PERFECT MONITORING

We consider the δ -discounted infinitely repeated prisoner's dilemma, the normalized stage-game payoffs of which are displayed below, where $g, l > 0$.

	C	D
C	1, 1	$-l, 1 + g$
D	$1 + g, -l$	0, 0

FIG. 1. Normalized Prisoners' Dilemma

In this section, monitoring is assumed perfect. Consider the family of behavior strategies defined as follows. Player i ($i = 1, 2$) plays C in stage 1, and in any subsequent stage t , his mixed action depends only on the outcome in stage $t - 1$. Denote by π_{a_i, a_j}^i the probability with which player i plays C conditional on the outcome (a_i, a_j) occurring in the previous period. Note that the behavior in period $t > 1$ is independent of t . For consistency with the following section, we will use lower-case subscripts, e.g. π_{cd}^j . Below is a diagram of the two-state machine that plays this strategy.⁶

We will show that for any payoff pair (v_1, v_2) in the square $V = (0, 1] \times (0, 1]$, there are strategies of this form that constitute an equilibrium of the repeated game (assuming sufficient patience) and obtain average payoffs (v_1, v_2) . In the next section, we show that equilibria of this form are robust to private monitoring. Obviously V is not the entire feasible, individually-rational set. The construction will make it clear that more cannot be achieved using strategies of this simple form. In section 4 we show how to

⁶Kandori and Obara [10] independently applied strategies of this form to repeated games with public monitoring.

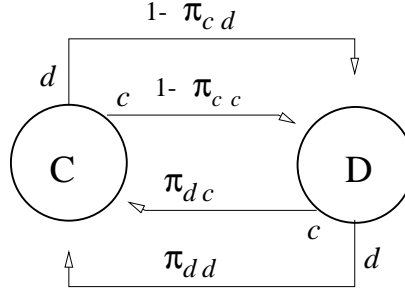


FIG. 2. Machine representation of π^i . To avoid clutter, we have left out the arrows that return to the preceding state.

modify these strategies to obtain any feasible individually rational payoff in a robust equilibrium. Let V° denote the interior of V .

Fix values V_C^i and V_D^i in $(0, 1]$, with $V_C^i > V_D^i$ for $i = 1, 2$, and let $\bar{\delta} < 1$ satisfy $\min_i \{V_C^i - V_D^i\} > 1 - \bar{\delta}$. The goal is to show that there exists $\bar{\delta}$ close enough to 1, such that for every $\delta \in (\bar{\delta}, 1)$ and for every $\hat{V}_C^i \in [V_C^i, 1]$, we can construct probabilities π^j such that (1) in any period in which j is playing C , player i is indifferent between actions C and D and obtains continuation value \hat{V}_C^i and (2), in any period in which player j is playing D , player i is indifferent between actions C and D and obtains continuation value V_D^i . From this it will follow that each player is indifferent between C and D after *every* history (a fact that is essential for the robustness argument), and player i gets average payoff \hat{V}_C^i , for $i = 1, 2$. The following four equations express these conditions:

$$\hat{V}_C^i = (1 - \delta) + \delta [\pi_{cc}^j \hat{V}_C^i + (1 - \pi_{cc}^j) V_D^i] \quad (1)$$

$$= (1 - \delta)(1 + g) + \delta [\pi_{cd}^j \hat{V}_C^i + (1 - \pi_{cd}^j) V_D^i] \quad (2)$$

$$V_D^i = -l(1 - \delta) + \delta [\pi_{dc}^j \hat{V}_C^i + (1 - \pi_{dc}^j) V_D^i] \quad (3)$$

$$= \delta [\pi_{dd}^j \hat{V}_C^i + (1 - \pi_{dd}^j) V_D^i] \quad (4)$$

Equation (1) is equivalent to the following

$$\pi_{cc}^j = \frac{\hat{V}_C^i - (1 - \delta) - \delta V_D^i}{\delta(\hat{V}_C^i - V_D^i)}$$

Because $\hat{V}_C^i - \delta V_D^i > V_C^i - V_D^i > 1 - \bar{\delta} > 1 - \delta$, the numerator and hence the fraction is greater than zero. It is no greater than one iff $\hat{V}_C^i - (1 - \delta) \leq \delta \hat{V}_C^i$

which is equivalent to $\hat{V}_C^i \leq 1$ which was assumed. For future reference, note that π_{cc}^j can be made interior if $\hat{V}_C^i < 1$.

Combining (1) and (2) and manipulating, we obtain

$$\pi_{cc}^j - \pi_{cd}^j = \left(\frac{1-\delta}{\delta} \right) \left(\frac{g}{\hat{V}_C^i - V_D^i} \right)$$

Since $\hat{V}_C^i - V_D^i \geq V_C^i - V_D^i > 0$,

$$0 < \pi_{cc}^j - \pi_{cd}^j \leq \left(\frac{1-\delta}{\delta} \right) \left(\frac{g}{V_C^i - V_D^i} \right)$$

and the upper bound can be made arbitrarily small by choosing δ sufficiently close to 1. Thus, $\pi_{cd}^j \in (0, 1)$ for sufficiently large δ .

We have shown that for each $\hat{V}_C^i \in [V_C^i, 1]$ there exist probabilities $\pi_{cc}^j \in (0, 1)$ and $\pi_{cd}^j \in (0, 1)$ such that if V_D^i is the value to player i when his opponent is playing D , then \hat{V}_C^i is the value to i when his opponent is playing C , and i is indifferent between C and D after such histories.

Now equation (4) reduces to

$$\pi_{dd}^j = \left(\frac{1-\delta}{\delta} \right) \left(\frac{V_D^i}{\hat{V}_C^i - V_D^i} \right)$$

Just as before, this is always positive and is bounded above by the expression obtained by substituting V_C^i for \hat{V}_C^i . Hence $\pi_{dd}^j \in (0, 1)$ for δ close enough to 1. Finally equations (3) and (4) combine to yield

$$\pi_{dc}^j - \pi_{dd}^j = \left(\frac{1-\delta}{\delta} \right) \left(\frac{l}{\hat{V}_C^i - V_D^i} \right)$$

which by the same argument makes π_{dc}^j an interior probability once δ is sufficiently large for each $\hat{V}_C^i \in [V_C^i, 1]$. Thus, interior probabilities can be found to ensure that when player j is playing D , player i is indifferent between C and D and obtains a continuation value V_D^i .

We conclude that for δ sufficiently close to 1 and for any value $\hat{V}_C^i \in [V_C^i, 1]$, there exists a strategy for player j satisfying equations 1-4. By the symmetric argument, there exists a strategy for player i which satisfies the analogous equations for player j . Furthermore, since these equations imply that each player is indifferent between C and D after every history, each strategy is a best-response to the other after every history. Thus, the

strategies form a subgame-perfect equilibrium of the perfect monitoring game and since each strategy is assumed to play C in the first stage, player i obtains average payoff $v_i = \hat{V}_C^i$.

3. A LIMIT FOLK THEOREM

The focus now turns to the repeated prisoners' dilemma with private monitoring. In this setting, players do not directly observe the actions chosen by their opponent. Following each stage of play, each player observes his own chosen action and a private signal which depends on the outcome in that stage. Let Σ_i be a finite set of signals for player i , and $\Sigma = \times_i \Sigma_i$. Assume that each Σ_i has at least two elements.

A monitoring technology is a collection $\{m(\cdot|a) : a \in \{C, D\}^2\}$ of probability distributions over signal profiles, one for each possible stage-game outcome. The marginal distribution over player i 's signal will be denoted $m^i(\cdot|a)$.

A monitoring m technology is ε -perfect if for each player i there is a set $c^i \subset \Sigma_i$ such that if $a_j = C$ then $m^i(c^i|a) > 1 - \varepsilon$ and if $a_j = D$ then $m^i(c^i|a) < \varepsilon$. In this paper, perfect monitoring means any ε -perfect monitoring technology for $\varepsilon = 0$.

Note that because our definition is stated in terms of marginal distributions alone, it entails no requirement on the correlation in a ε -perfect monitoring technology. In particular, it includes as special cases the two polar cases often studied in the literature: imperfect public monitoring (in which case the players' signal sets coincide and the joint distribution exhibits perfect correlation), and conditionally independent monitoring (in which case $m^i(\cdot|a)$ and $m^j(\cdot|a)$ are independent distributions for each action profile a .)

We restrict attention to ε -perfect monitoring technologies and write $m_{a_i, a_j}^i = m^i(c^i|a_i, a_j)$, so for example m_{DC}^i represents the probability that i observes a signal in c^i when i plays D and j plays C . Let d^i denote the complement in Σ_i of c^i .

The δ -discounted repeated prisoners' dilemma with private monitoring technology m will be denoted $G^\infty(\delta, m)$. A strategy π^i in $G^\infty(\delta, m)$ specifies a mixed action for each history of own-actions and observed signals. We will restrict attention to strategies which depend only on histories of length 1 and condition only on the player's own action and the events c^i and d^i . By analogy to the previous section, write e.g. π_{dc}^i for the probability with which player i plays C following a stage in which i played D and observed a signal in c^i . Set $\pi_\emptyset^i = 1$.

For any two strategy profiles $\pi, \tilde{\pi}$ of this form, define the distance $|\pi - \tilde{\pi}|$ to be $\max |\pi_s^i - \tilde{\pi}_s^i|$ where the maximum is over players $i = 1, 2$ and states s of the machine.

DEFINITION 3.1. A subgame-perfect equilibrium π of $G^\infty(\delta, m_0)$ is **robust to private monitoring** if for every $\epsilon > 0$ there exists $\varepsilon > 0$ such that for all ε -perfect monitoring technologies m , there is a sequential equilibrium $\tilde{\pi}$ of $G^\infty(\delta, m)$, within ϵ distance of π and with payoffs within ϵ of the payoffs under π .

A sufficient condition for a pair (π^1, π^2) to be a sequential equilibrium of $G^\infty(\delta, m)$ with payoffs (v_1, v_2) is for the mixtures to solve the following equations for $V_C^i = v_i, i = 1, 2, j \neq i$.

$$\begin{aligned} V_C^i &= (1 - \delta) + \delta V_C^i \left[m_{CC}^j \pi_{cc}^j + (1 - m_{CC}^j) \pi_{cd}^j \right] \\ &\quad + \delta V_D^i \left[(1 - m_{CC}^j)(1 - \pi_{cd}^j) + m_{CC}^j(1 - \pi_{cc}^j) \right] \end{aligned} \quad (5)$$

$$\begin{aligned} V_C^i &= (1 - \delta)(1 + g) + \delta V_C^i \left[m_{CD}^j \pi_{cc}^j + (1 - m_{CD}^j) \pi_{cd}^j \right] \\ &\quad + \delta V_D^i \left[(1 - m_{CD}^j)(1 - \pi_{cd}^j) + m_{CD}^j(1 - \pi_{cc}^j) \right] \end{aligned} \quad (6)$$

$$\begin{aligned} V_D^i &= -l(1 - \delta) + \delta V_C^i \left[m_{DC}^j \pi_{dc}^j + (1 - m_{DC}^j) \pi_{dd}^j \right] \\ &\quad + \delta V_D^i \left[(1 - m_{DC}^j)(1 - \pi_{dd}^j) + m_{DC}^j(1 - \pi_{dc}^j) \right] \end{aligned} \quad (7)$$

$$\begin{aligned} V_D^i &= \delta V_C^i \left[m_{DD}^j \pi_{dc}^j + (1 - m_{DD}^j) \pi_{dd}^j \right] \\ &\quad + \delta V_D^i \left[(1 - m_{DD}^j)(1 - \pi_{dd}^j) + m_{DD}^j(1 - \pi_{dc}^j) \right] \end{aligned} \quad (8)$$

Equations (5) and (6) state that player i gets average payoff V_C^i when he plays either C or D in any stage in which player j plays C . Equations (7) and (8) state that player i gets average payoff V_D^i when he plays either C or D in any stage in which player j plays D . If these equations are satisfied, then player i is always indifferent between his two actions and is therefore willing to play any mixed strategy after every history.

When monitoring is perfect, $m_{*,D}^j = 1 - m_{*,C}^j = 0$ and these equations reduce to equations 1-4 from the previous section. There it was shown that provided $V_C^i \in (0, 1]$ and $V_D^i < V_C^i$, a solution exists. Moreover, a solution $\tilde{\pi}$ exists for every $\bar{V}_C^i \in [V_C^i, 1]$, and the mixtures $\tilde{\pi}_{cd}^j$ and $\tilde{\pi}_{dc}^j$ can be chosen to be interior. We can now fix $\tilde{\pi}_{cc}^j$ and $\tilde{\pi}_{dd}^j$ at the solution and solve for

V_C^i, V_D^i, π_{cd}^j , and π_{dc}^j :

$$\begin{aligned}\pi_{dc}^j &= \frac{\delta \bar{\pi}_{dd}^j \mu_D (\mu_C - g(1 - m_{CC}^j)) + l(1 - \delta(\bar{\pi}_{cc}^j - \bar{\pi}_{dd}^j) + \bar{\pi}_{dd}^j m_{DD}^j)}{\delta g \mu_D (m_{CC}^j - 1) + \mu_C (\mu_D - l m_{DD}^j)} \\ \pi_{cd}^j &= \frac{\delta \mu_C \bar{\pi}_{cc}^j (\mu_D + l \bar{\pi}_{dd}^j) + g \mu_D (\delta (\bar{\pi}_{cc}^j m_{CC}^j - \bar{\pi}_{dd}^j) - 1)}{\delta g \mu_D (m_{CC}^j - 1) + \mu_C (\mu_D - l m_{DD}^j)} \\ V_C^i &= \bar{V}_C^i + \frac{g \mu_D (1 - m_{CC}^j) (1 + \delta (\bar{\pi}_{dd}^j - 1)) + \delta \mu_C l (\bar{\pi}_{cc}^j - 1) m_{DD}^j}{\mu_C \mu_D (\delta (\bar{\pi}_{cc}^j - \bar{\pi}_{dd}^j) - 1)} \\ V_D^i &= \bar{V}_D^i + \frac{\delta [g \mu_D \bar{\pi}_{dd}^j (1 - m_{CC}^j) + l \mu_C \bar{\pi}_{cc}^j m_{DD}^j] - l \mu_C m_{DD}^j}{\mu_C \mu_D (\delta (\bar{\pi}_{cc}^j - \bar{\pi}_{dd}^j) - 1)}\end{aligned}$$

where we have written $\mu_D = m_{DC}^j - m_{DD}^j$ and $\mu_C = m_{CC}^j - m_{CD}^j$.

These equations define the left-hand side variables as continuous functions of the monitoring parameters. Since $\bar{\pi}_{cd}^j$ and $\bar{\pi}_{dc}^j$ were interior, it follows that for $m_{*,D}^j$ and $1 - m_{*,C}^j$ sufficiently close to zero, there exists a solution (π^j, V_D^i, V_C^i) to equations 5-8 with mixtures π^j and values V_C^i, V_D^i arbitrarily close to $(\bar{\pi}^j, \bar{V}_C^i, \bar{V}_D^i)$. Player j therefore has a strategy against which every strategy of player i is a best-response and achieves payoff V_C^i . This establishes the following theorem.

THEOREM 3.1 (Robust folk theorem). *Let $(v_1, v_2) \in V$. There exist $\bar{\delta} \in (0, 1)$ such that for all $\delta \in (\bar{\delta}, 1)$, and for all $(\hat{v}_1, \hat{v}_2) \in V$ such that $\hat{v}_i \geq v_i$, there exists a robust subgame perfect equilibrium of $G^\infty(\delta, m_0)$ with payoffs (\hat{v}_1, \hat{v}_2) .*

Now suppose $\bar{V}_C^i \in (0, 1)$ and recall that this implies that $\bar{\pi}_{cc}^j \in (0, 1)$. We can now fix $(\bar{V}_C^i, \bar{V}_D^i) \in (0, 1)^2$ and solve equations 5-8 for the following explicit relations between the equilibrium strategies and the monitoring probabilities.

$$\begin{aligned}\pi_{cc}^j &= \frac{(\bar{V}_C^i - \delta \bar{V}_D^i)(m_{CC}^j - m_{CD}^j) + (1 - \delta)(g + m_{CD}^j - m_{CC}^j(1 + g))}{\delta (m_{CC}^j - m_{CD}^j) (\bar{V}_C^i - \bar{V}_D^i)} \\ \pi_{cd}^j &= \frac{(\bar{V}_C^i - \delta \bar{V}_D^i)(m_{CC}^j - m_{CD}^j) + (1 - \delta)(m_{CD}^j - (1 + g)m_{CC}^j)}{\delta (m_{CC}^j - m_{CD}^j) (\bar{V}_C^i - \bar{V}_D^i)} \\ \pi_{dc}^j &= \frac{(1 - \delta) [l(1 - m_{DD}^j) - \bar{V}_D^i(m_{DD}^j - m_{DC}^j)]}{\delta (m_{DC}^j - m_{DD}^j) (\bar{V}_C^i - \bar{V}_D^i)} \\ \pi_{dd}^j &= \frac{(1 - \delta) [(m_{DC}^j - m_{DD}^j)\bar{V}_D^i - lm_{DD}^j]}{\delta (m_{DC}^j - m_{DD}^j) (\bar{V}_C^i - \bar{V}_D^i)}\end{aligned}$$

One can easily verify that these define equilibrium behavior strategies as continuous functions of the monitoring probabilities m^j . Since all mixtures in π^j are interior, it follows that for $m_{*,D}^j$ and $1 - m_{*,C}^j$ sufficiently close to zero, equations 5-8 can be solved for probabilities π^j , and the symmetric set of equations can be solved for probabilities π^i . These solutions are sequential equilibria of the private monitoring game with values $(\bar{V}_C^1, \bar{V}_C^2)$. This proves the limit folk theorem

THEOREM 3.2 (Limit folk theorem). *Let $(v_1, v_2) \in V^\circ$. There exist $\bar{\delta}, \varepsilon \in (0, 1)$ such that for all $\delta \in (\bar{\delta}, 1)$ and for all ε -perfect monitoring technologies m , there exists a sequential equilibrium of $G^\infty(\delta, m)$ with payoffs (v_1, v_2) .*

The order of quantification was not proven above. In the previous section, it was shown that if there is a solution for some δ then there is a solution for any larger $\hat{\delta}$. One can easily verify that the same holds here: if there are probabilities that solve equations 5-8 for some ε and δ , then there are probabilities that solve the equations for ε and any $\hat{\delta} > \delta$.

4. EXTENDING TO THE FULL PAYOFF SET

We have shown that any payoff pair in V can be obtained in a robust equilibrium. This leaves out much of the feasible, individually rational set of payoffs. In particular, the only efficient payoff pair in V is the symmetric payoff $(1, 1)$. Unfortunately, nothing outside of V can be obtained robustly using strategies of the simple form considered above. To see this, note

that in our equilibrium, player i is indifferent between C and D after every history. This implies that his long-run average payoff must be equal to the payoff he would get by playing C after every history. Obviously such a payoff cannot exceed 1.

However, more can be achieved using more complicated (yet still finite-state) strategies. We sketch the idea here, and prove it formally below. Return to the perfect monitoring case. Fix $V_D^1 > 0$ close to zero and $V_C^2 = 1$. Let V_C^1 be any value in $(V_D^1, 1]$. We have shown in theorem 3.1 that for δ close enough to 1, there exist robust equilibrium strategies (π^1, π^2) which obtain values (V_C^1, V_C^2) , and player 2's strategy can be chosen so that player 1 obtains value V_D^1 whenever player 2 plays D .

Consider the following strategies. Player 1 plays C in stage 1, then proceeds with π^1 starting in stage 2. Player 2 plays D in stage 1. In stage 2, player 2 switches to strategy π^2 but starts in state D if he observes that player 1 has played D in stage 1. If 2 observes that 1 has played C , 2 starts in state C .

Let $(v_1, v_2) = (1 - \delta)(-l, 1 + g) + \delta(V_C^1, V_C^2)$. The above strategies form an equilibrium with payoffs (v_1, v_2) if the following incentive constraints are satisfied.

$$(1 - \delta)(1 + g) + \delta V_C^2 > (1 - \delta) + \delta V_C^2 \quad (9)$$

$$(1 - \delta)(-l) + \delta V_C^1 > \delta V_D^1 \quad (10)$$

These hold for sufficiently large δ . Note that the payoffs (v_1, v_2) are outside of V . Now we claim that this equilibrium is robust to private monitoring. To see this, note that in stage two, each player is indifferent between his two actions regardless of the outcome in stage 1. Each is therefore willing to behave as the strategy demands even after erroneous signals. The continuation payoffs V_C^1, V_C^2, V_D^2 are continuous in the monitoring perturbation because the continuation strategies are robust. Thus, because the first-stage incentive constraints are satisfied with strict inequalities, they will continue to hold for small perturbations of m_0 . We show below that by increasing the number of stages in which (C, D) is played, any feasible individually rational payoff vector can be sustained in a robust equilibrium.

To prove this formally, we make the simplifying assumption that $g - l \leq 1$ so that mutual cooperation is not Pareto-dominated by any other profile. Under this assumption any individually rational payoff by a convex combination of the set V and a single additional point. The theorem is true without the assumption, and could be proved by the same line of argument using strategies that alternate between (C, D) and (D, C) .

Let V^* be the interior of the set of feasible, individually rational payoff vectors.

THEOREM 4.1. *For each $v \in V^*$ there is a $\bar{\delta} \in (0, 1)$ such that for all $\delta \in (\bar{\delta}, 1)$ there is a robust subgame perfect equilibrium of $G^\infty(\delta, m_0)$ with payoffs v .*

Proof. Let V_1^* be the subset of V^* in which player 1's payoff is at least 1 (and hence player 2's payoff is less than 1). We will prove the claim for all payoffs in V_1^* . Together with theorem 3.1, and the symmetric argument when the players roles are reversed, this will establish the result.

Under the assumption that $g - l \leq 1$, V_1^* is in the convex hull of $V^\circ \cup \{(1 + g, -l)\}$ (see figure 1). In particular, for each $v \in V_1^*$, there exists a w

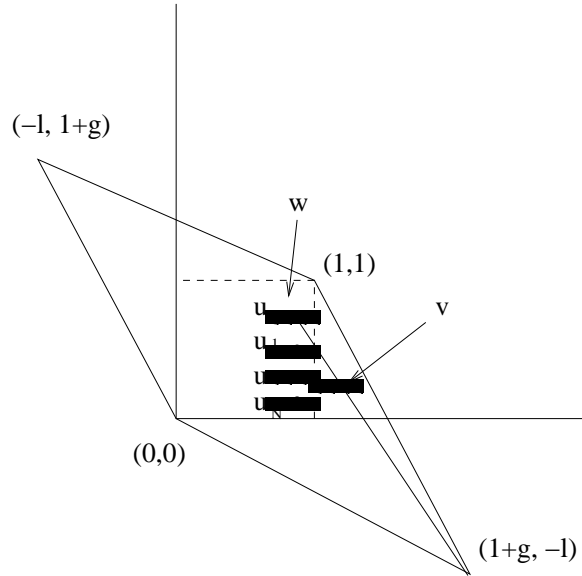


FIG. 3. Prisoner's Dilemma feasible payoff set.

in the interior of V and a $t \in (0, 1)$ such that

$$v = (1 - t)[1 + g, -l] + tw.$$

For every $\delta > t$ define $N(\delta) = \max\{N : \delta^N > t\}$. Note that $\lim_{\delta \rightarrow 1} \delta^{N(\delta)} = t$. Choose $w(\delta)$ to satisfy

$$v = (1 - \delta^{N(\delta)})[1 + g, -l] + \delta^{N(\delta)}w(\delta).$$

Since $\lim_{\delta \rightarrow 1} w(\delta) = w$, for every δ sufficiently close to 1, $w(\delta)$ is in the interior of V . For each δ , define the following sequence.

$$\begin{aligned} u_0 &= w_2(\delta) \\ u_{k+1} &= u_k - \frac{l(1-\delta)\delta^k}{t} \end{aligned}$$

The sequence is monotonically decreasing and

$$u_{N(\delta)} = u_0 - \frac{l(1-\delta^{N(\delta)})}{t}.$$

which as δ approaches 1, converges to

$$u^* = w_2 - l \left(\frac{1-t}{t} \right) \quad (11)$$

By the definition of t , $t = \frac{v_2+l}{l+w_2}$. Substituting into (11):

$$\begin{aligned} u^* &= w_2 - l \left(\frac{w_2 - v_2}{v_2 + l} \right) \\ &> v_2 \end{aligned}$$

which is strictly positive since v is in the interior of V .

Since we have shown that $\lim_{\delta \rightarrow 1} u_{N(\delta)} > 0$, for all δ sufficiently close to 1, the entire sequence $\{u_k\}_{k=1}^{N(\delta)}$ is included in the interval $(0, 1)$. Now by Theorem 3.1 there is a $\bar{\delta}$ such that for every $\delta \in (\bar{\delta}, 1)$, every pair (w_1, u_k) for $k = 1, \dots, N(\delta)$ is a robust equilibrium payoff.

For any such δ , we now construct a robust subgame perfect equilibrium which achieves payoffs v . Player 1 will play D for each of the first $N = N(\delta)$ stages independent of history. Beginning in stage $N + 1$, player 1 will use a continuation strategy that depends on the number of signals in d^1 he has observed. In particular, if m signals in d^1 were observed, then player 1 will play a strategy which makes 2 indifferent among all strategies and which gives 2 a continuation payoff of u_m .

Player 2 will play C in each of the first N stages, regardless of history, and in stage $N + 1$ will begin playing a strategy which makes player 1 indifferent among all strategies and which gives player 1 a continuation payoff of w_1 .

Since player 2's behavior is independent of any history through stage N , it is clearly optimal for player 1 to play his dominant action in each of these stages, whatever the history. Furthermore, player 2's continuation strategy beginning in stage $N + 1$ makes player 1 indifferent among all strategies, so player 1 is willing to condition his play at that stage as specified.

Since every one of player 1's continuation strategies at stage $N+1$ makes player 2 indifferent among all strategies, player 2's specified continuation strategy at that stage is optimal. It remains only to show that player 2 is playing optimally at all histories prior to that stage.

Player 2's continuation payoffs and hence optimal continuation strategy depends only on player 1's continuation strategy. This continuation strategy depends only on the history observed by 1. We will show that whatever the history observed by 1, every one-stage deviation by player 2 is *strictly* sub-optimal.

Suppose $k < N$ stages have been played and $m \leq k$ signals in d^1 have been observed by player 1. A one-shot deviation by 2 gains l today but increases the number of d^1 signals from m to $m+1$ and so reduces the continuation payoff beginning in stage $N+1$ by $(\frac{1-\delta}{t})\delta^m l$. The net discounted gain is

$$(1 - \delta)l \left[1 - \frac{\delta^{N+m-k}}{t} \right]$$

which is strictly negative since $m \leq k$ and therefore $\delta^{N+m-k} > t$.

We have shown that the strategies form a subgame perfect equilibrium. When monitoring is perfect, the profile (D, C) will be played with certainty in the first N stages and therefore the equilibrium payoff vector is $(1 - \delta^N)[1 + g, -l] + \delta^N w(\delta) = v$.

Finally, we demonstrate that the equilibrium is robust. To begin with, since each of the continuation payoffs in stage $N+1$ is a robust payoff, for all sufficiently small perturbations of perfect monitoring, there exist sequential equilibria which achieve payoffs as close as necessary to the original payoffs. In particular, we will choose the perturbation small enough so that the original behavior in the first N stages remains strictly optimal.

As before, since player 2's behavior is independent of any signals observed in the first N stages, player 1 will continue to optimally play D in each of the first N stages independent of history. Also as before, since player 1's play in the first N stages is independent of history, player 2's optimal continuation strategy at any stage $k \leq N$ depends only on the number of d^1 signals that 1 has observed. Under perfect monitoring, regardless of the number of d^1 signals already received, it was strictly optimal for 2 to play C . By choosing the perturbation small enough we can make the distribution over future signals as well as the continuation values beginning in stage $N+1$ as close as necessary to their perfect monitoring counterparts in order to preserve the strict optimality at each of the (finitely many) histories through stage N . ■

The proof of Theorem 4.1 shows that every payoff v in $V^* \setminus V$ and every discount factor δ sufficiently close to 1, there is a sequential equilibrium whose payoffs are arbitrarily close to v provided monitoring is close enough

to perfect. Since in the strategies employed in the proof, the number of stages in which (D, C) played grows to infinity as δ goes to 1, it may appear that the upper bound on the size of the perturbation shrinks as δ approaches 1. However, this is not the case and in fact an application of the Ellison [6] trick shows that if v can be achieved with a given monitoring technology for discount factor δ , then payoffs arbitrarily close to v can be achieved under the same monitoring technology for any discount factor sufficiently close to 1.

Since the argument is now standard in the literature, we only sketch it here. In the equilibrium in the proof of Theorem 4.1, for discount factor δ , each player has a strictly optimal action in each of the first N stages. We can therefore find a small interval $[\delta, \delta + z]$ of higher discount factors in which these initial actions remain optimal. (Play beginning in stage N will have to change slightly to maintain indifference and the original continuation payoffs. Theorem 3.1 ensures this is possible.) Since the payoffs are continuous in the discount factor, we can choose z small enough to ensure that all payoffs are within an arbitrarily small neighborhood of v .

Now for every discount factor $\hat{\delta}$ close enough to 1, there is an integer M such that $\hat{\delta}^M \in [\delta, \delta + z]$. For any such discount factor, we can construct a strategy which treats the repeated game as M component games played in stages $1, M + 1, \dots$ and $2, M + 2, \dots$, etc. In these games, the effective discount factor is δ^M and therefore play in each component game is an equilibrium achieving payoffs arbitrarily close to v and the same is true of the overall strategy profile.

5. EXTENSIONS

In this section, we outline briefly how the method for finding robust subgame perfect equilibria can be extended to more general games.

5.1. Two-Player Games

Let $G = (A_1, A_2, g_1, g_2)$ be a finite two-player normal form game, where A_i are the action sets and g_i payoff functions. Let $G^\infty(\delta)$ be the infinitely repeated game with G as the stage game and δ as the discount factor. Consider two strategy profiles (a'_1, a'_2) and (a''_1, a''_2) and assume that the following condition on the payoff functions is satisfied: For $i \in \{1, 2\}$, there are \underline{v}_i and $\bar{v}_i > \underline{v}_i$ such that

$$i) \max_{a_i \in A_i} g_i(a_i, a''_{-i}) < \underline{v}_i \text{ and}$$

$$ii) \min \{g_i(a'_i, a'_{-i}), g_i(a''_i, a'_{-i})\} > \bar{v}_i.$$

The first step is again the construction of a subgame perfect equilibrium in the game with perfect monitoring where the players mix between (a'_i, a''_i) . In analogy to the previous sections, let $\pi_{a_i, a_j}^i \in \Delta(A_i)$ be the probability distribution on player i 's pure actions conditional on observing outcome (a_i, a_j) at the previous stage, and let $V_{a_j}^i$ be as before.

THEOREM 5.1. *For any $v \in (\underline{v}_1, \bar{v}_1) \times (\underline{v}_2, \bar{v}_2)$, there is a $\bar{\delta} < 1$ such that v can be supported as SPE payoff in $G^\infty(\delta)$ for all $\delta > \bar{\delta}$ using strategies where $\pi_{a_i, a_j}^i(a'_i) > 0$, $\pi_{a_i, a_j}^i(a''_i) > 0$ and $\pi_{a_i, a_j}^i(\tilde{a}_i) = 0$ for $\tilde{a}_i \notin \{a'_i, a''_i\}$.*

Proof. See appendix. ■

Using this theorem, it is easy to see that in some games, payoffs below those resulting from the unique dominant strategy equilibrium of the stage game are sustainable in robust subgame perfect equilibrium and hence in games with small imperfections in the monitoring technologies. Consider e.g. figure 4.

	a_1	a_2
a_1	1, 1	-1, 0
a_2	0, -1	-2, -2

FIG. 4. Payoffs below the dominant strategy equilibrium can be supported.

The theorem above implies that any $v \in (-1, 0)^2$ can be supported as a robust equilibrium payoff. Observe also that the pure strategy Nash equilibrium payoff cannot be approximated by a mixed strategy equilibrium of the type as described above.⁷

Two corollaries follow immediately from the theorem. Let $e = (e_1, e_2)$ be a pure strategy Nash equilibrium profile of the stage game. Denote the stage game equilibrium payoffs by v_1^e and v_2^e respectively.

COROLLARY 5.1. *If there is an $a = (a_1, a_2)$ such that for $i = 1, 2$*

$$\min \{g_i(e_i, a_{-i}), g_i(a_i, a_{-i})\} > v_i^e$$

then the set of robust subgame perfect equilibrium payoffs has a non empty interior. Furthermore, we can require that player i uses at each stage either a_i or e_i .

⁷Although such payoffs can be obtained by strategies which play (a_1, a_1) for sufficiently many periods before reverting to an equilibrium in the set $(-1, 0)^2$. These strategies are robust by the argument made in section 3.

As an example of this corollary, consider the discretized Cournot competition model with linear demands. For concreteness, let $A_i = \{0, 1, 2, \dots, 20\}$ and $g_i(a_i, a_j) = a_i(20 - a_i - a_j)$. A Nash equilibrium of the stage game is $e_i = e_j = 7$. The monopoly price is supported by action profile $a_1 = a_2 = 5$. Observe that $g_i(e) = 42$, $g_i(a) = 50$ and $g_i(e_i, a_j) = 56$. As a result, we conclude that any $v \in (42, 50)^2$ can be supported as a robust subgame perfect equilibrium for sufficiently high δ .

For the second corollary, let $m_i \in A_i$ be the action that minmaxes player j . Let v_i denote the pure strategy minmax payoff of each player.

COROLLARY 5.2. *If there is an (a_1, a_2) such that $\min\{g_i(m_i, a_{-i}), g_i(a_i, a_{-i})\} > v_i$ for $i \in \{1, 2\}$, then the set of robust subgame perfect equilibrium payoffs has a non empty interior. Furthermore, we can require that player i uses at each stage either a_i or m_i .*

5.2. N-Player Case

We conclude this section by analyzing an example of a symmetric N -player game. The basic difficulty in comparison to the two player case is that the transitions between the states of the machines equivalent to those described in the previous section are no longer linear in the randomizations of the other players (even after conditioning on own actions). As a result, the probabilities cannot be solved for as before by simple linear algebra. In this subsection, we show that for δ close to unity, we can recover approximate linearity and show the robustness by an application of the implicit function theorem.⁸

Consider the following version of the N -player Prisoner's Dilemma. $A_i = \{C, D\}$ for $i = 1, \dots, N$, and $g_i(D, a_{-i}) = n_i(a_{-i}) = g_i(C, a_{-i}) + 1$ where $n_i(a_{-i})$ is the number of players different from i that play C in profile a_{-i} . Let $G^\infty(\delta)$ stand for the infinitely repeated version of this game with discount factor δ .

Suppose that all players are using strategies of the following form:

$$\pi_{a_i, n_i}^i = \Pr\{C \mid a\}.$$

In words, each player bases her own randomizations solely on her own past actions and the number of other players that cooperated, not their identity. The next theorem shows that outcomes arbitrarily close to the efficient outcome can be supported

⁸Bhaskar and Obara [2] also analyze the N -player prisoner's dilemma. Using an extension of the Sekiguchi (1997) approach, they obtain an approximately efficient symmetric equilibrium.

THEOREM 5.2. *For any $v \in (0, N - 2)$, there is a $\bar{\delta}$ such that for all $\delta > \bar{\delta}$, $G^\infty(\delta)$ has a completely mixed SPE in which each player obtains a payoff of v .*

Proof. See appendix. ■

In order to obtain the result on games with almost perfect monitoring, fix a δ for which there is a completely mixed SPE with payoff vector v . We can make the same continuity argument as in the previous section to conclude that whenever monitoring is close enough to perfect, a sequential equilibrium with payoff vector v exists.

6. CONCLUSION

There are many unresolved questions in the context of discounted repeated games with private monitoring. In particular, unlike the public monitoring case, there is no known folk theorem for fixed monitoring technologies.⁹ In this note, we have followed the literature and sought after the weaker result: a limit folk theorem for nearly perfect monitoring. An important goal for future research is to characterize equilibrium payoffs when monitoring is private but far from perfect. The strategies we have used in this paper can sustain some cooperation under less than perfect monitoring technologies and equilibrium can be characterized by an analogous system of dynamic programming equations. The payoffs that can be supported are thus those values for which the system can be solved for probabilities. Determining the full potential of this approach is ongoing research.

7. APPENDIX

Proof of Theorem 5.1

Pick $V_{a_j^i}$ and $V_{a_j^{i'}}$ such that $\underline{v}_i < V_{a_j^{i'}} < V_{a_j^i} < \bar{v}_i$. Suppose that j randomizes between $a_j^{i'}$ and a_j^i in each period so that $\pi_{a_j, a_i}^j(\tilde{a}_j) = 0$ for $\tilde{a}_j \notin \{a_j^i, a_j^{i'}\}$. The claim is proved if there are probabilities $\pi_{a_j, a_i}^j(a_j^{i'})$ such that the following conditions hold:

$$\begin{aligned} V_{a_j^{i'}} &= (1 - \delta) g_i(a_i^{i'}, a_j^{i'}) + \delta \left[\pi_{a_j, a_i}^j(a_j^{i'}) V_{a_j^{i'}} + \left(1 - \pi_{a_j, a_i}^j(a_j^{i'})\right) V_{a_j^i} \right] \\ &= (1 - \delta) g_i(a_i^i, a_j^{i'}) + \delta \left[\pi_{a_j, a_i}^j(a_j^{i'}) V_{a_j^{i'}} + \left(1 - \pi_{a_j, a_i}^j(a_j^{i'})\right) V_{a_j^i} \right]. \end{aligned}$$

⁹Lehrer [11] and [12] prove folk theorems for non-discounting criteria

For all $a_i \notin \{a'_i, a''_i\}$,

$$V_{a'_j}^i \geq (1 - \delta) g_i(a_i, a'_j) + \delta \left[\pi_{a'_j, a_i}^j(a''_j) V_{a'_j}^i + \left(1 - \pi_{a'_j, a_i}^j(a''_j)\right) V_{a'_j}^i \right].$$

$$\begin{aligned} V_{a'_j}^i &= (1 - \delta) g_i(a''_i, a'_j) + \delta \left[\pi_{a'_j, a'_i}^j(a''_j) V_{a'_j}^i + \left(1 - \pi_{a'_j, a'_i}^j(a''_j)\right) V_{a'_j}^i \right] \\ &= (1 - \delta) g_i(a'_i, a'_j) + \delta \left[\pi_{a'_j, a'_i}^j(a''_j) V_{a'_j}^i + \left(1 - \pi_{a'_j, a'_i}^j(a''_j)\right) V_{a'_j}^i \right], \end{aligned}$$

and also for all $a_i \notin \{a'_i, a''_i\}$,

$$V_{a'_j}^i \geq (1 - \delta) g_i(a_i, a'_j) + \delta \left[\pi_{a'_j, a_i}^j(a''_j) V_{a'_j}^i + \left(1 - \pi_{a'_j, a_i}^j(a''_j)\right) V_{a'_j}^i \right]$$

Our condition on payoffs ensures that $g_i(a_i, a'_j) < V_{a'_j}^i$ for all $a_i \in A_i$, and $g_i(a''_i, a'_j) > V_{a'_j}^i$, and $g_i(a'_i, a'_j) > V_{a'_j}^i$.

The first set of equalities and inequalities can, in fact, be solved as a set of equalities, and we get:

$$\pi_{a'_j, a_i}^j(a''_j) = \frac{V_{a'_j}^i - (1 - \delta) g_i(a_i, a'_j) - \delta V_{a'_j}^i}{\delta (V_{a'_j}^i - V_{a'_j}^i)}.$$

Since $V_{a'_j}^i > V_{a'_j}^i$ and $g_i(a_i, a'_j) < V_{a'_j}^i$, $0 < \pi_{a'_j, a_i}^j(a''_j) < 1$ whenever δ is sufficiently close to 1.

The second set of equalities yields:

$$\begin{aligned} \pi_{a'_j, a'_i}^j(a''_j) &= \frac{(1 - \delta) (V_{a'_j}^i - g_i(a'_i, a'_j))}{\delta (V_{a'_j}^i - V_{a'_j}^i)} \\ \pi_{a'_j, a'_i}^j(a''_j) &= \frac{(1 - \delta) (V_{a'_j}^i - g_i(a'_i, a'_j))}{\delta (V_{a'_j}^i - V_{a'_j}^i)}. \end{aligned}$$

Again, $\pi_{a'_j, a'_i}^j(a''_j)$ and $\pi_{a'_j, a'_i}^j(a''_j)$ are probabilities for δ sufficiently close to 1.

The second set cannot be satisfied as equalities, in general. To see this solve for the probabilities:

$$\pi_{a'_j, a_i}^j(a''_j) = \frac{(1 - \delta) V_{a'_j}^i - (1 - \delta) g_i(a_i, a'_j)}{\delta (V_{a'_j}^i - V_{a'_j}^i)}.$$

If $g_i(a_i, a'_j) < V_{a'_j}^i$, and $V_{a''_j}^i < V_{a'_j}^i$, the solution is negative.

This does not cause any problems, however, since we can pick an arbitrary \bar{V}^i with $V_{a''_j}^i < \bar{V}^i < V_{a'_j}^i$ and probabilities $\pi_{a'_j, a_i}^j(a''_j)$ that yield for all $a_i \notin \{a'_i, a''_i\}$:

$$\bar{V}^i = (1 - \delta) g_i(a_i, a'_j) + \delta \left[\pi_{a'_j, a_i}^j(a''_j) V_{a''_j}^i + \left(1 - \pi_{a'_j, a_i}^j(a''_j)\right) V_{a'_j}^i \right].$$

A similar set of conditions holds for player j , and given that each player is indifferent after each history between actions a'_i and a''_i , and given that all other actions yield a weakly lower payoff, the randomizations above are best responses.

Proof of Theorem 5.2

Construct the following sequence of numbers:

$$V(0) = v_0, \quad V(n) = v_0 + \frac{n}{(N-1)}(v - v_0) \quad \text{for } n \in \{1, \dots, N-1\}. \quad (12)$$

We show that each $V(n)$ can be interpreted as the value to a player in the game when n other players cooperate at the current stage.

To simplify notation, write the strategies as:

$$\pi_{C, n_i}^i = 1 - \gamma_n \quad \text{and} \quad \pi_{D, n_i}^i = \beta_n.$$

Note that we have dropped the superscript as we will now restrict attention to symmetric strategies. We are interested in the existence of equilibria in strategies of this type for δ large enough. The dynamic programming equations become then for $n \in \{0, \dots, N-1\}$:

$$\begin{aligned} V(n) &= (1 - \delta)n + \delta E_{D, n} V(n'), \\ V(n) &= (1 - \delta)(n - 1) + \delta E_{C, n} V(n'), \end{aligned} \quad (13)$$

where $E_{C, n} V(n')$ denotes the expected value from tomorrow on conditional on the action profile today. The transition probabilities are nonlinear in (γ_n, β_n) since they are obtained from two binomial distributions.

Suppose that the strategies used by the players depend on δ as follows:

$$\gamma_n(\delta) |_{\delta=1} = 0, \quad \beta_n(\delta) |_{\delta=1} = 0 \quad \text{for all } n \in \{0, \dots, N-1\}.$$

In this case, the equations in (13) are trivially satisfied at $\delta = 1$, and in particular, the sequence given in (12) solves the system.

Write $\gamma(\delta) = (\gamma_0(\delta), \dots, \gamma_{N-1}(\delta))$ and $\beta(\delta) = (\beta_0(\delta), \dots, \beta_{N-1}(\delta))$. We want to use the implicit function theorem to conclude that the nonlinear

system, (13) has a solution $(\gamma(\delta), \beta(\delta))$ for δ in a neighborhood of 1. Next, to guarantee that the solutions are indeed probabilities, we need to show that for all n , $\frac{\partial \gamma_n}{\partial \delta}(1) < 0$, $\frac{\partial \beta_n}{\partial \delta}(1) < 0$.

To apply the implicit function theorem, we treat (13) as a system of $2N$ equations that implicitly define the $2N$ probabilities $(\gamma(\delta), \beta(\delta))$. Write this system as

$$V(\gamma, \beta; \delta) = 0,$$

with γ, β as the endogenous variables and δ as the exogenous variable.

Viewing (13) as a system on N pairs of equations, denote by V_{nD} the first function of the n th pair and V_{nC} the second. We now evaluate the partial derivatives of the system at the original solution (for $\delta = 1$) with respect to the endogenous variables. For $1 \leq n \leq N - 1$,

$$\frac{\partial V_{nD}}{\partial \gamma_{n-1}} = \frac{\partial V_{nC}}{\partial \gamma_n} = n(V(n-1) - V(n)),$$

and for $0 \leq n \leq N - 2$,

$$\frac{\partial V_{nD}}{\partial \beta_n} = \frac{\partial V_{nC}}{\partial \beta_{n+1}} = (N - n - 1)(V(n+1) - V(n)).$$

It can easily be verified that this system has full rank. The implicit function theorem then guarantees the existence of a differentiable solution mapping $(\gamma(\delta), \beta(\delta))$ in the neighborhood of $\delta = 1$. To check the sign of the derivatives, we apply the chain rule. Note

$$D_\delta V(\gamma, \beta; \delta) = \begin{pmatrix} V(0) \\ V(0) + 1 \\ \vdots \\ V(k) - k \\ V(k) - k + 1 \\ \vdots \\ V(N-1) - (N-1) \\ V(N-1) - (N-2) \end{pmatrix}$$

Observe that the system of equations,

$$D_{\gamma, \beta} V(\gamma, \beta, \delta) D_\delta(\gamma, \beta) = -D_\delta V(\gamma, \beta; \delta),$$

can be solved in blocks of size 2. This makes the comparative statics relatively easy. For each $n = 1, \dots, N - 3$,

$$\begin{aligned} & \begin{pmatrix} n(V(n-1) - V(n)) & (N-n-1)(V(n+1) - V(n)) \\ (n+1)(V(n) - (V(n+1))) & (N-n-2)(V(n+2) - V(n+1)) \end{pmatrix} \begin{pmatrix} \frac{\partial \gamma_n}{\partial \delta} \\ \frac{\partial \beta_{n+1}}{\partial \delta} \end{pmatrix} \\ &= - \begin{pmatrix} V(n) - n + 1 \\ V(n+1) - n - 1 \end{pmatrix}. \end{aligned}$$

Using the fact that $V(n+1) - V(n) = V(k) - V(k-1)$ for all $k, n \in \{0, 1, \dots, N-1\}$, it is easy to see that the derivatives have the right sign. Hence for δ close enough to one, the system can be solved for probabilities. This implies that for such discount factors, the game has a completely mixed SPE with a payoff v .

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