

DISCUSSION PAPER NO. 48

DESIGN OF AN OPTIMIZATION SYSTEM FOR  
UNIVERSITY USE

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# "Design of an Optimization System for University Use"

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## ABSTRACT

An optimization system for university use is quite different from MPSX or OPTIMA. This paper deals with the design and development of a Multi Purpose Optimization System (MPOS) now in use on Northwestern University's CDC 6400, which has evolved in the following stages:

- a) the "stand-alone" program library - which includes over twenty linear, integer and non-linear programming algorithms,
- b) the "batch" MP system - developed to permit the user to state LP/QP/IP problems in English and standard MP notation. The system, coded in FORTRAN, consists of a language pre-processor and individual overlays corresponding to the various solution algorithms,
- c) the "online" MP system - designed to be used interactively at a teletype.

Our design objectives have been to stress ease of use, expandability and interactive optimization. The system's modular design allows for incorporation of new or improved algorithms. They, in turn, can be tested with an archive of test problems. Usage statistics are collected to better plan for future software development or improvement.

The advantages of interactive optimization are indicated: maximize the utility in real time between the user and the computational process, facilitate the learning of basic principles (e.g., pivoting, cuts), provide options to interrupt computation and permit the user to interpose decisions based on his experience and intuition and to spend more time on the modelling aspects of optimization.

## 1. Introduction

The essential role of a university computer center is to provide the university community the maximum opportunity for making use of the computer as a problem-solving or information-processing tool. A person with an optimization problem is especially eager to have dependable computer codes requiring a minimum of concern for system details outside the scope of his primary interest.

Most mathematical programming (MP) systems, such as IBM's MPSX, CDC's OPTIMA/OPHELIE/APEX, or UNIVAC's UMPIRE, are directed at the solution of very large problems stemming from corporate or industry models. An optimization system for university use is quite different from MPSX or OPTIMA. First, the software supplied by computer manufacturers is proprietary and very expensive. Second, the majority of university users have small or medium scale problems. Third, the typical user is not inclined to solve optimization problems on the computer if he has to spend a great deal of time learning about the architecture of MP systems. More often than not, their documentation is not designed for a novice user.

Our design objectives have been to stress ease of use, expandability, and interactive optimization. Our library of MP algorithms has evolved from a batch version to an online version in a little more than two years. Our system was developed to permit the user to state his problem in English and standard MP notation, to access any available linear programming (LP), integer programming (IP), or quadratic programming (QP) codes, and to edit/store/retrieve his problem files.

## 2. Motivation

Our motivation arose from practical needs:

- We wanted to provide a simple to use and economical LP/IP/QP system to some 500 management, engineering and economics students

enrolled in introductory optimization courses. The fact that about 150 management students are enrolled in the evening program, prompted the development of an online version which would enable them to solve optimization problems "at their office desks" if they have a terminal available.

- The choice between CDC's OPTIMA and standard "canned" programs was not very appealing. On one hand, CDC's OPTIMA is expensive to use in the classroom because of its high set-up cost. On the other hand, the standard canned programs differ in the way they accept input data and therefore force the student to follow different instructions for each algorithm. Hence, the need for a language pre-processor to create the various tableaux required by the different algorithms and act as a common input language.

- Finally, we realized that there are many pedagogical advantages in man-machine interaction. The learning process of the computational aspects of optimization is accelerated. Also, an interactive optimization system is economical in real-time. The student needs a few runs to get an answer to his problem and can, therefore, spend more time on the modelling aspects of optimization.

### 3. An overview of optimization software at Northwestern University

The field of mathematical software is an emerging scientific discipline and interactive applied mathematical systems have been developed in parallel with new on-line operating systems, see e.g. [25], [33]. The most notable efforts in MP systems design and implementation of the batch processing type have been those of Beale [4], Dantzig [9], [5], and Orchard-Hays [40]. For a recent survey report on the computational aspects of MP see White [41]. In this section, we briefly describe our optimization software.

### 3.1 The "stand-alone" library

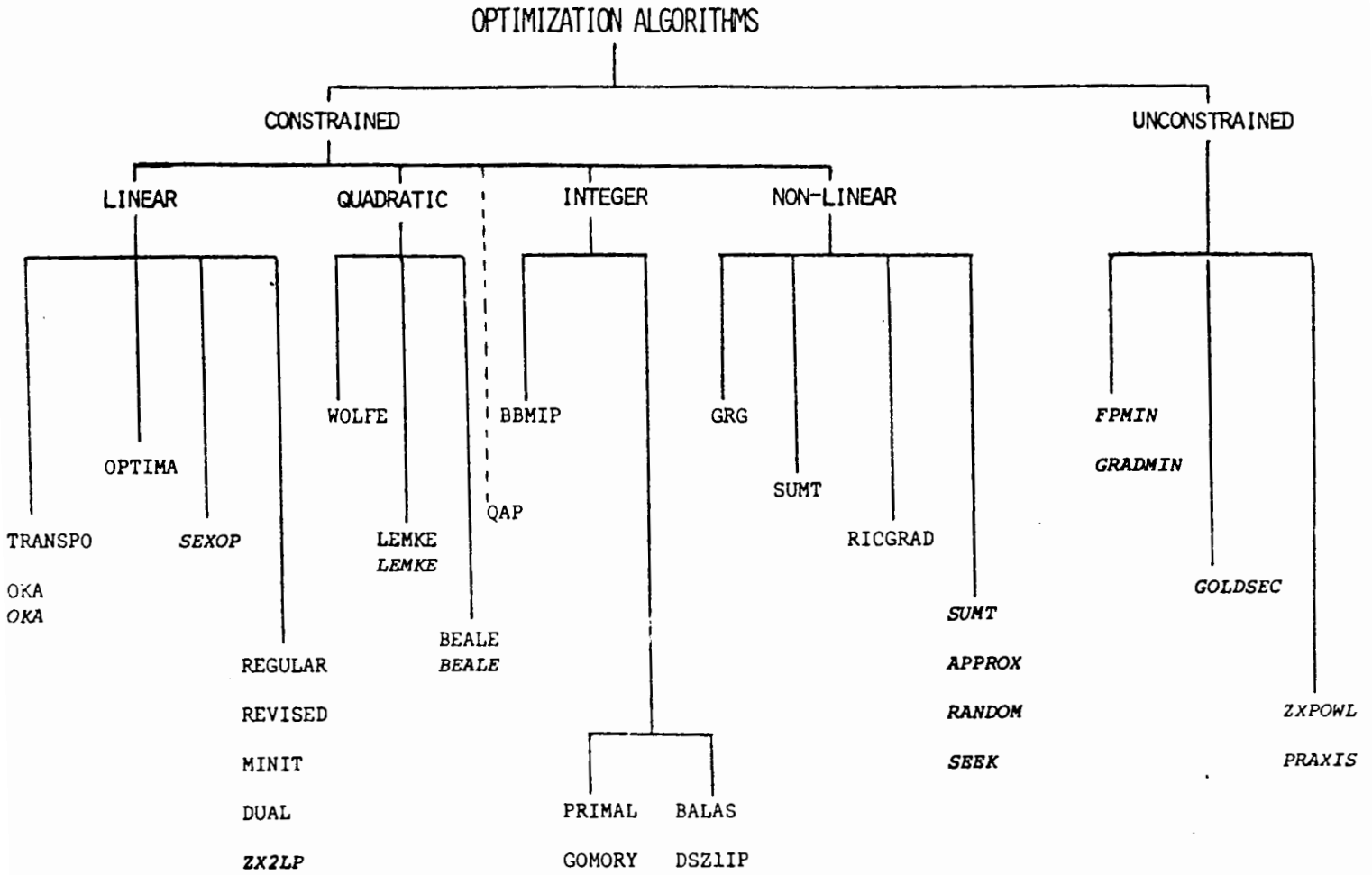
This library of "canned" programs comprises many algorithms of constrained and unconstrained optimization and is best described in Tables 1 and 2. These programs have been collected from several sources: some were developed locally, some were contributed by interested users, some were obtained from research laboratories or other university computer centers. We check them out and document them individually. The stand-alone library is tested using an archive of standard test problems (Colville [7] , Haldi [20] , Himmelblau [22] , Trauth and Woolsey [37] ) or problems contributed by our own users.

### 3.2 The MPOS system (Multi Purpose Optimization System)

This system was developed to permit the user to state his problem in English and standard MP notation. The user has access to the LP/IP/QP algorithms. For example, he may input the following:

WOLFE	algorithm name
MAXIMIZE	
X1 + X2 -0.5X1*X1 + X1*X2 - X2*X2	objective function definition
CONSTRAINTS	
X1 + X2 .LE. 2	constraints definition
2X1 + 3X2 .GE. 6	
PRINT 1	print each iteration
OPTIMIZE	
STOP	stop computation

The system consists of a language pre-processor and individual overlays corresponding to the various solution algorithms. With the exception of a few assembly language subroutines, the MP system is coded in FORTRAN. The system's modular design allows for incorporation of new or improved algorithms.



(Subroutines in italics)

TABLE I

OPTIMIZATION ALGORITHMS

---

Constrained optimization

Linear:

TRANSPO	Transportation algorithm (primal-dual)	[8]
OKA/ <i>OKA</i>	Out-of-kilter algorithm	[8]
OPTIMA	CDC's LP system	[30]
SEXOP	Subroutines for Experimental Optimization (decomposition and GUB constraints)	[29]
REGULAR	2-phase simplex algorithm	[8]
REVISED	Revised simplex algorithm	[8]
MINIT	Primal-dual algorithm	[34]
DUAL	Dual simplex algorithm	[8]
ZX2LP	Revised simplex algorithm	[24]

Quadratic:

WOLFE	Wolfe's quadratic simplex algorithm	[26]
LEMKE/ <i>LEMKE</i>	Lemke's complementary pivot algorithm	[26]
BEALE/ <i>Beale</i>	Beale's quadratic programming algorithm	[26]

Integer:

BBMIP	Branch and bound mixed integer prog.	[36]
PRIMAL GOMORY	Glover/Young/Harris primal cutting plane Gomory's dual cutting plane algorithm	[15,21,39] [16]
BALAS DSZ1IP	Balas/Glover additive algorithm Lemke and Spielberg direct search alg.	[2,14] [27]
QAP	Graves and Whinston quadratic assignment	[17]

Non-linear:

SUMT	Fiacco and McCormick SUMT algorithm	[10]
GRG	Abadie and Carpentier generalized reduced gradient algorithm	[1]
RICGRAD	Greenstadt's ricochet gradient algorithm	[19]
<i>SEEK</i>	Hooke and Jeeves direct search alg.	[23]
<i>APPROX</i>	Griffith and Stewart linear approximation algorithm	[18]
<i>RANDOM</i>	Random search subroutine	[38]

Unconstrained optimization

<i>FPMIN</i>	Fletcher and Powell descent method	[11]
<i>GRADMIN</i>	Fletcher and Reeves conjugate gradient	[12]
<i>GOLDSEC</i>	Golden section search	[38]
<i>ZXPOWL</i>	Powell's method without derivatives	[31]
<i>PRAXIS</i>	Brent's method without derivatives	[6]

### 3.3 The interactive MPOS system

This system is an extension of the batch system and is designed to be used interactively at a terminal. It has editing and file manipulation capabilities, such as, generating files to be processed in batch mode when output is voluminous. The interactive system can be used in two computation modes. The "direct" mode simulates the batch version. The "assisted" mode interacts with the user (see § 5).

Detailed statistics on the use of these programs permit us to better plan for future software development or improvement.

### 4. MPOS from the designer's view-point

Northwestern University's Vogelback Computing Center has the following hardware/software configuration:

- CDC 6400, 65K (60 bit) words central memory, 250K Extended Core Storage, SCOPE 3.3 Operating System and NU's own ONLINE time-sharing system.

The need for modularity resulted in adopting an overlay program structure for MPOS:

- The main (0,0) overlay selects the algorithm and determines if computation mode is batch or online.
- The (1,0) primary overlay is the language pre-processor. It decodes statements such as

5.3X1 - 2X1\*X1 + X2\*X1     or

X1 - 2X2    .LE. 4            or

BOUNDS

X1    .LE.    10

X2    .LE.    5



and performs syntax checking. Its main function is to produce a data file of  $a_{ij}$ ,  $b_i$ ,  $c_j$ , etc. for input to the different algorithms or to CDC's OPTIMA.

- The other primary overlays are the algorithm overlays.

Each algorithm overlay computes the exact amount of central memory required by the problem size. This dynamic storage allocation permits the loading of the system in less than 13000 decimal words. It also allows for stacking problems until the STOP command is encountered. .

In all, MPOS includes some 60 subroutines. All but two are in FORTRAN. A full description of the system can be found in the system documentation manual.

#### 5. MPOS from the user's view-point

We have tried to keep the number of control cards to a minimum, and make the input look like standard MP notation. All the user has to do is become familiar with a number of keyword commands (see Table 3).

\*\*\*\*\* PROBLEM NUMBER 22 \*\*\*\*\*

Example 2

```

MAXIMIZE
X1+X10
CONSTRAINTS
X1+X2-3 = 5
-- FATAL ERROR NUMBER 4
   MISPLACED SIGN OR NUMERIC CONSTANT

X1+3.5X2.LE. X5+1
-- FATAL ERROR NUMBER 3
   VARIABLE ON RIGHT HAND SIDE OF RELATIONAL

*DECK TEST21
-----
-- PROBLEM ABORTED DUE TO FATAL ERRORS.

```

Example 3

```

EDITOR READY
? TEXT. ←————— Enter text of QP problem
? MAXIMIZE
? X1 + X2 -0.5X1X1 - X1X2 +X1X2
? CONSTRAINTS
? X1+X2 <3
? 2X1+3X2 > 6
? OPTIMIZE
? ZE ←————— Terminate text
LAST LINE 60
? LIST. ←————— Produce listing of text
10 MAXIMIZE
20 X1 + X2 -0.5X1X1 - X1X2 +X1X2
30 CONSTRAINTS
40 X1+X2 <3
50 2X1+3X2 > 6
60 OPTIMIZE
? PEPLACE/+X1/-X2/20 ←————— Correct '+X1' to '-X2' in line 20
20 X1 + X2 -0.5X1X1 - X1X2 -X2X2
? AL,GOQP ←————— Execute problem

```

```

--- QP ---
VERSION 1.5
WOLFE'S QUADRATIC SIMPLEX

```

PROBLEM NO. 1

## INITIAL TABLEAU

BASIC VAR.	CP	VALUES
9	-1.	3.0000
10	-1.	6.0000
11	-1.	1.0000
12	-1.	1.0000

Z(J)-C(J)\* -11.000

----- FINAL SOLUTION AT ITERATION 4

BASIC VAR.	CB	VALUES
3	0.	.40000
8	0.	.80000
1	0.	1.8000
2	0.	.80000

← Summary results

Z(J)-C(J)\* 0.

-1.100000 = VALUE OF OBJECTIVE FUNCTION

.0230 CP SECONDS FOR 4 ITERATIONS

SUCCESSFUL EXECUTION

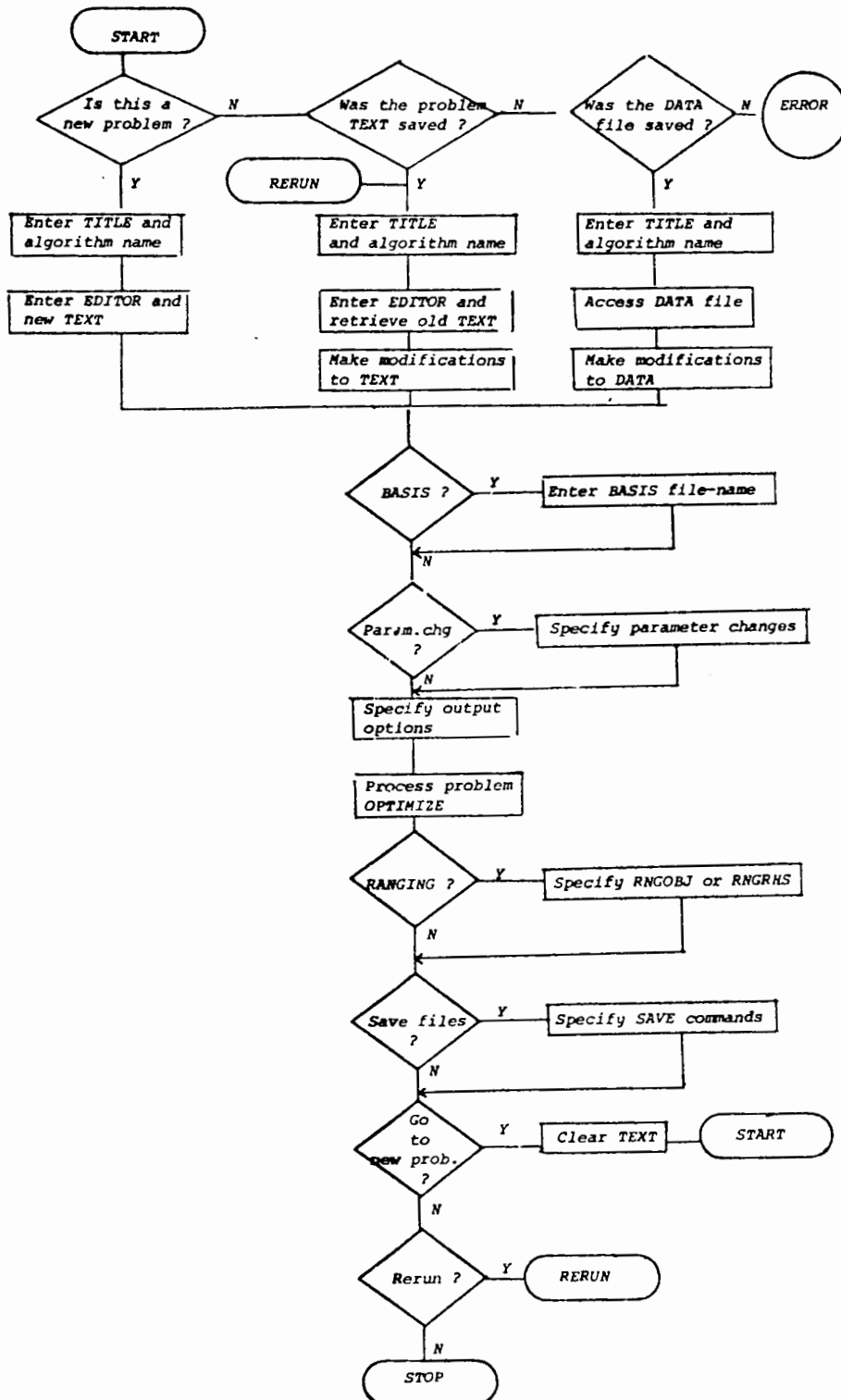
? NEW\*CPFILE ←————— Save problem text on file QPPROB

? LEAF

<i>TITLE</i>	Any alphanumeric title
<i>REVISED...WOLFE...BBMIP</i>	Name of the optimization algorithm
<i>INTEGER VARIABLES</i>	List of integer variables List of variable names
<i>MAXIMIZE MINIMIZE</i>	Objective function definition
<i>CONSTRAINTS CONSTRAINTS n</i>	Constraints definition
<i>BOUNDS BNDALL value BNDOBJ value</i>	Variable upper bounds Objective function upper-bound
<i>EPSILON value TOLERANCE value LIMIT n</i>	Perturbation factor Accuracy check Maximum number of iterations
<i>PRINT PRINT i</i>	Print initial and final tableaus Print every i-th tableau
<i>RNGOBJ RNGRHS</i>	Range on objective function coefficients Range on right-hand side coefficients
<i>SAVEFILE file-name GETFILE file-name READ file-name</i>	Store data file Retrieve data file Read alternate input file
<i>PACKED TABLEAU MATRIX</i>	Packed input Tableau input Matrix input
<i>FORMAT</i>	Variable format
<i>MODIFY</i>	Data file modifications
<i>BASIS file-name SAVEBASIS file-name</i>	Input basis Save basis
<i>OPTIMIZE END</i>	Initiate problem solution
<i>STOP</i>	Stop run
<i>*</i>	(any comment card)

TABLE 3

When solving problems interactively, the user has the option between the DIRECT and ASSISTED solution modes. In the former case, he is in fact simulating the batch version. In the latter case, he interacts with the system via questions and answers. Only summary output is printed locally. A number of examples follow the online flowchart for the ASSISTED mode.



\*\*\*\*\* PROBLEM NUMBER 9 \*\*\*\*\*

Example 1

```

MAXIMIZE
*****
*
* MAXIMIZE PROFIT FROM MANUFACTURE OF
*   TABLES, CHAIRS, DESKS, AND BOOKCASES
*
* VARIABLE DEFINITIONS,
*   X1 = TABLES
*   X2 = CHAIRS
*   X3 = DESKS
*   X4 = BOOKCASES
*
*****
12X1 + 5X2 + 15 X3 + 10X4
CONSTRAINTS
**
* CONSTRAINTS DUE TO SALES SCHEDULE
**
X3 .GE. 15
X4 .GE. 10
**
* CONSTRAINTS DUE TO AVAILABLE RESOURCES
**
* BOARD FEET OF TYPE I LUMBER
5X1+X2+9X3+12X4 .LE. 1500
* BOARD FEET OF TYPE II LUMBER
3X1+2X2+5X3+10X4 .LE. 1000
* MAN HOURS TO UTILIZE
2X1+3X2+4X3+5X4 = 800
OPTIMIZE
    
```

\*\*\*\*\*  
\* PROBLEM NUMBER 9 \*  
\*\*\*\*\*

ITERATION	1	IN VAR-	4	OUT VAR-	8	-WMIN=	-765.000
ITERATION	2	IN VAR-	3	CUT VAR-	7	-WMIN=	-690.000
ITERATION	3	IN VAR-	6	CUT VAR-	10	-WMIN=	-277.500
ITERATION	4	IN VAR-	2	CUT VAR-	11	-WMIN=	0.

ENTERING PHASE II

ITERATION	5	IN VAR-	1	CUT VAR-	5	-ZMIN=	3373.00
-----------	---	---------	---	----------	---	--------	---------

\*\*\*\*\*  
\* PROBLEM NUMBER 9 \*  
\*\*\*\*\*

SUMMARY OF RESULTS

VARIABLE NO.	VARIABLE NAME	BASIC NON-BASIC	ACTIVITY LEVEL	OPPORTUNITY COST
1	X1	B	219.0000000	--
2	X2	B	84.0000000	--
3	X3	B	15.0000000	--
4	X4	B	10.0000000	--
5	--SLACK	NB	--	3.8000000
6	--SLACK	NB	--	33.0000000
7	--ARTIF	NB	--	-3.8000000
8	--ARTIF	NB	--	-33.0000000
9	--SLACK	B	66.0000000	--
10	--SLACK	NB	--	5.2000000
11	--ARTIF	NB	--	-1.8000000

MAXIMUM VALUE OF THE OBJECTIVE FUNCTION = 3373.000000

CALCULATION TIME WAS .0300 SECONDS FOR 5 ITERATIONS.

Example 4

----- PROBLEM NUMBER 2 -----

ENTERING EDITOR  
EDITOR READY  
Cleared

24 LINES READ  
7 5. BMIP  
7 6. INTEGER LIST  
7 7. 0A 0B BA BC AC AK CK  
7

Now solve the same problem as an I.P.

Insert BMIP command  
Insert INTEGER list  
Insert list of INTER var.

Execute problem  
Branch and Bound Mixed Integer Programming

----- PROBLEM NUMBER 2 -----

NO. OF ROWS INCLUDING OBJ. FN. 7  
NO. OF COLS. INCLUDING RIGHT-HAND-SIDE 9  
NO. OF INT. VARIABLES 7

ITERATION 1: OUT ROW- 6, IN COL- 8, -Z=0.  
ITERATION 2: OUT ROW- 1, IN COL- 2, -Z= 11.000  
ITERATION 3: OUT ROW- 2, IN COL- 6, -Z= 22.000  
ITERATION 4: OUT ROW- 3, IN COL- 2, -Z= 33.000  
ITERATION 5: OUT ROW- 4, IN COL- 3, -Z= 44.000  
ITERATION 6: OUT ROW- 1, IN COL- 1, -Z= 47.000  
ITERATION 7: OUT ROW- 2, IN COL- 3, -Z= 55.000

----- PROBLEM NUMBER 1 -----

Begin new problem

ENTERING EDITOR  
EDITOR READY  
Cleared

7 \* MIN-COST NETWORK FLOW STATED AS A L.P. PROBLEM

7 VARIABLE LIST  
7 0A 0B BA BC AC AK CK K0  
7 \* SOURCE NODE = 0, SINK NODE = K, OTHER NODES=A,B,C

7 MINIMIZE  
7 40A + 0B + 2BA + 3BC + 6AC + AK + 2CK  
7 \* CONSERVATION OF FLOWS

7 K0 - 0A - 0B = 0  
7 AK + CK - K0 = 0  
7 0A + BA - AC - AK = 0  
7 0B - BA - BC = 0  
7 BC + AC - CK = 0

7 K0=11

7 \* CAPACITATED LINKS

7 \* BOUNDS  
7 0A<10  
7 0B<8  
7 BA<5  
7 BC<10  
7 AC<2  
7 AK<7  
7 CK<4

7 \* OPTIMIZE

7 \* TERMINATE TEXT  
Begin execution of problem

LAST LINE 240  
\* MIN-COST NETWORK FLOW STATED AS A L.P. PROBLEM

USING REGULAR  
Regular simplex is used by default

----- PROBLEM NUMBER 2 -----

CONTINUOUS SOLUTION ---

OBJECTIVE FUNCTION= 55.00000 AT ITERATION 7. TIME= .052

SUMMARY OF RESULTS

VARIABLE NO.	NAME	BASIC/NON-B	INTEG/CONTIN	ACTIVITY LEVEL	OPPORTUNITY COST
1	0A	B	I	3.0000000	
4	BC	B	I	4.0000000	
7	CK	B	I	4.0000000	
3	BA	B	I	4.0000000	
0	--ARTIF	B	C	-0.0000000	
8	K0	NB	I	11.0000000	
2	0B	NB	I	8.0000000	
5	AC	NB	I		-1.0000000
6	AK	NB	I	7.0000000	5.0000000
					-2.0000000

TOLERANCE SET AT 60.50000 AT ITERATION 7

CONTINUOUS SOLUTION IS INTEGER SOLUTION.

OPTIMALITY ESTABLISHED AT ITERATION 7

-----END OF PROBLEM. TOTAL TIME WAS .095 SECS.

PROBLEM COMPLETE, TYPE STOP TO TERMINATE RUN, OR EDIT TO MODIFY PROBLEM.  
7 STOP

----- PROBLEM NUMBER 1 -----

SUMMARY OF RESULTS

VARIABLE NO.	VARIABLE NAME	BASIC/NON-BASIC	ACTIVITY LEVEL	OPPORTUNITY COST
1	0A	B	3.0000000	
2	0B	B	8.0000000	
3	BA	B	4.0000000	
4	BC	B	4.0000000	
5	AC	NB	7.0000000	
6	AK	B	4.0000000	
7	CK	B	11.0000000	
8	K0	NB	0.0000000	
9	--ARTIF	B	0.0000000	
10	--ARTIF	NB	3.0000000	
11	--ARTIF	NB	5.0000000	
12	--ARTIF	NB	2.0000000	
13	--ARTIF	NB	-7.0000000	
14	--ARTIF	NB	1.0000000	
15	--SLACK	B	1.0000000	
16	--SLACK	NB	1.0000000	
17	--SLACK	B	6.0000000	
18	--SLACK	B	2.0000000	
19	--SLACK	B	0.0000000	
20	--SLACK	NB		
21	--SLACK	B		

MINIMUM VALUE OF THE OBJECTIVE FUNCTION = 55.000000

CALCULATION TIME WAS .0220 SECONDS FOR 7 ITERATIONS.

PROBLEM COMPLETE, TYPE STOP TO TERMINATE RUN, OR EDIT TO MODIFY PROBLEM.  
7 EDIT

## 6. Evaluation of the system and conclusions

In the previous sections we described a software system designed to solve linear, integer and quadratic optimization problems using well established algorithms. Because of its simple structure and repertoire of algorithms, MPOS has become one of the most widely used applications programs at Northwestern University. As a result, students feel less "intimidated" by non-linear optimization models and they find it practical to experiment with extensions of LP models. Our emphasis on modularity facilitates expandability. The system is flexible enough to incorporate new algorithms. A student or researcher can "plug" into our system a new algorithm module without much concern for input/output and evaluate it with our archive of test problems.

Our system has proven to be economical from the view-points of usage and maintenance. The user needs few runs to get a solution to his problem. If, for example, an integer programming algorithm fails, the user can switch to another one by changing one card or editing a line. Because the interactive system needs a very small amount of central memory (5000 words) to load, response times are very good. Maintenance of the system is greatly simplified by the overlay structure which lends itself to "decentralized programming" activities. Good internal documentation and machine independence (of the batch version) facilitate transportability.

It is quite evident that our problem sizes are limited by the capacities of central memory and Extended Core Storage. Our system is not designed to compete with the new large-scale MP systems, such as those under development at Stanford's System Optimization Laboratory or MIT's NBER/Computer Research Center in Economics and Management Science. Users with large problems turn to APEX I and II, a commercial LP/MIP system we just obtained from CDC.

The reliability of the algorithms depends on the efficiency of the numerical methods used to implement them. For example, to speed computation time we perform row elimination procedures in assembly language. We are also studying the adoption of new matrix decomposition techniques (Bartels and Golub [3], Saunders [35]) to improve numerical accuracy.

Too often, system designers have concentrated on effective utilization of hardware and software resources with little attention paid to man/machine dialogue. Stated simply, these systems have been designed from the inside, out. To be user-effective, systems must be designed from the outside, in, because computer users come from a spectrum of individuals with different levels of knowledge and different kinds of problems. The value of this MP system will be judged by the users of MPOS and by the amount by which MPOS reduces the total time it takes a person to solve an optimization problem.

This system represents the first step toward our goal of a general interactive optimization system designed to "graft" the computer to students or researchers involved in optimization modelling. Existing algorithms are highly formalized and do not allow much human intervention and interaction. We believe that an interactive optimization system fulfills a great need in both research and teaching. Faced with a LP problem with multiple objective functions, the decision maker learns to recognize good solutions and the relative importance of (competing) objectives through sensitivity analysis performed online. In fact, phases of computation alternate with phases of decision.

Although the present system excludes NLP algorithms, there are instances where the user may change solution strategy through the use of different algorithms. For example, starting with a cutting-plane method and terminating with a branch and bound. There are a number of practical and theoretical problems in non-linear optimization where mixed-strategy can lead to new insights: convergence properties of algorithms, interaction of algorithms and criteria for switching from one to



the other, solution of non-convex or ill-structured problems, best matching of an algorithm to a given class of problems. As more experience is gained with our existing system, we will have a better perspective for designing the interactive non-linear optimization system.

In closing, we can state that this operational system has pedagogical advantages, facilitates man-machine problem-solving in that it brings the union of optimization and computing. It is responsive to the user needs and background and can be used as an inexpensive tool for training users in the power and pitfalls of various optimization techniques.

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